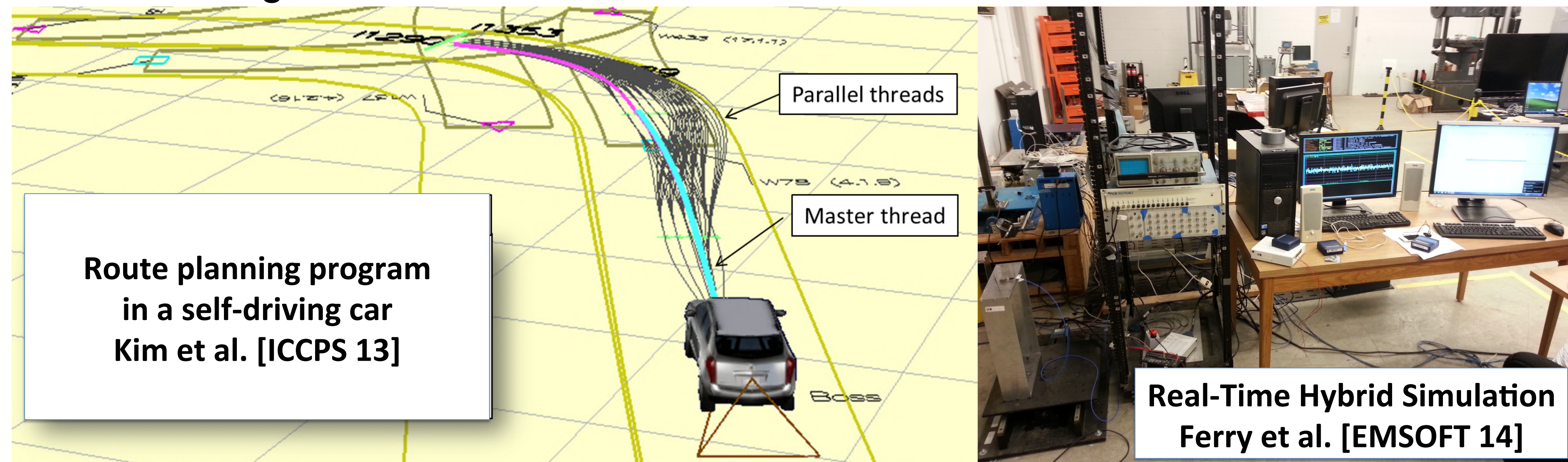


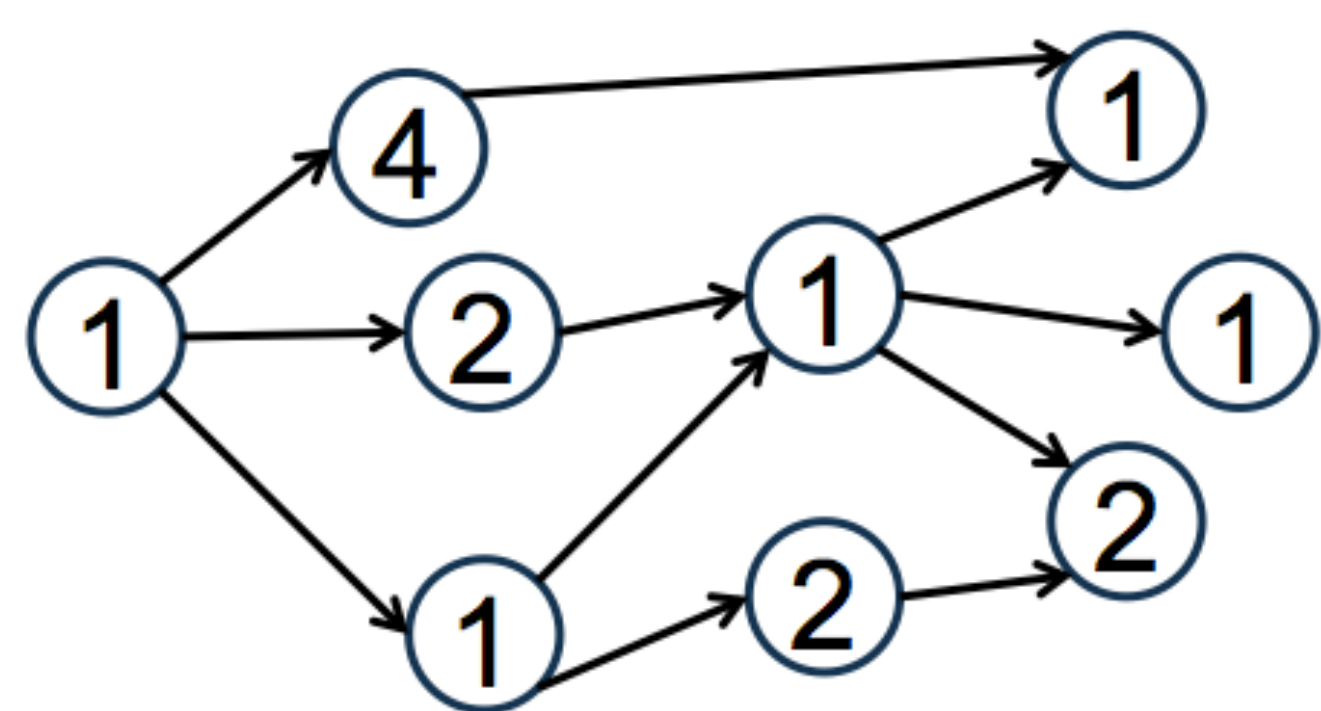
XPS: Real-Time Scheduling of Parallel Tasks

Motivation: Real-Time Systems Are Everywhere

- Real-time tasks are tasks with *deadlines*.
- They arise when computation interacts with physical world.
- Modern real-time applications are more computationally intensive and have tighter deadlines.



Goal: Parallel Real-Time Systems



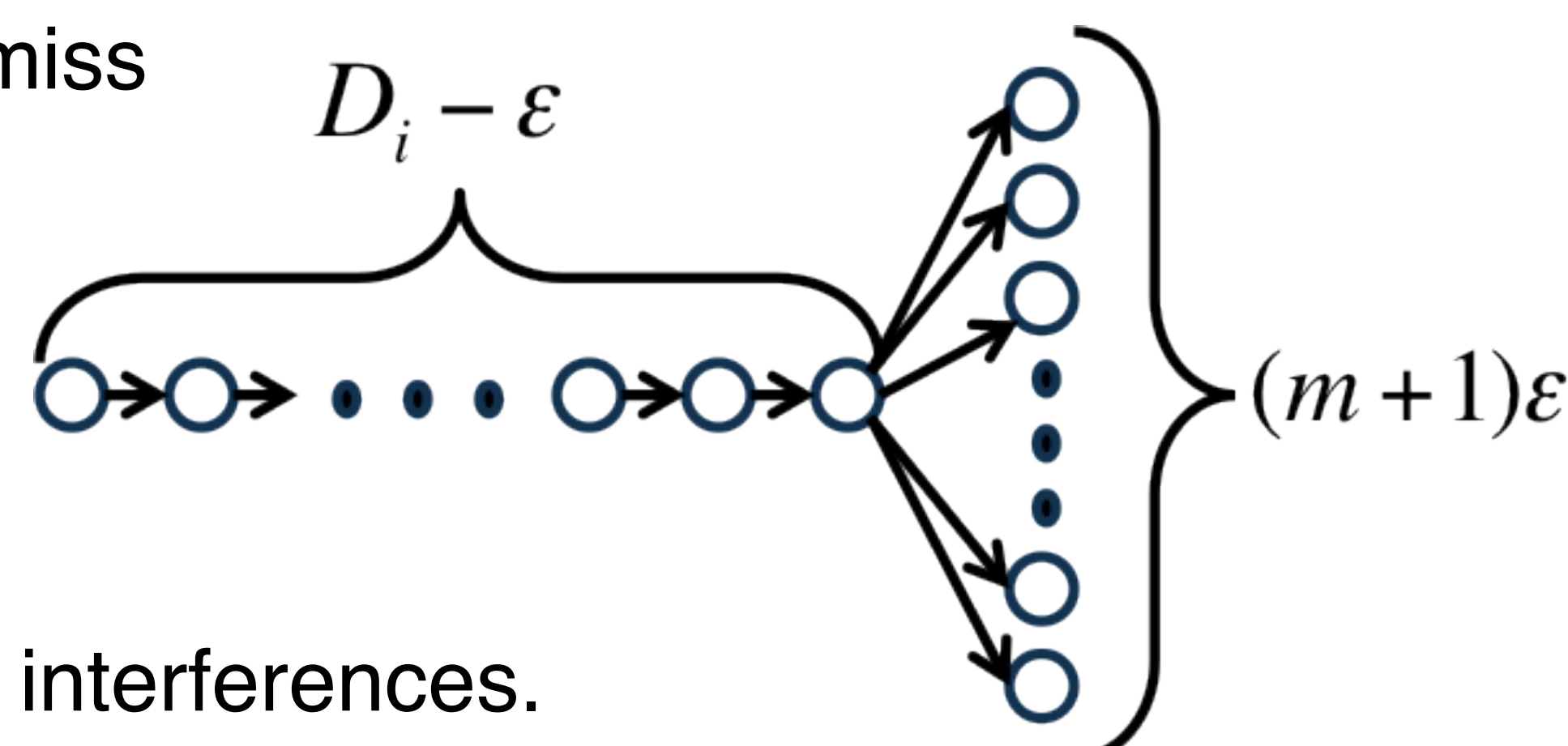
Execution time C_i
Critical-path length L_i
Utilization $u_i = C_i/D_i$

- Develop theoretical foundations and system implementations to support parallel real-time tasks.
- Task model: Directed Acyclic Graph (DAG).

- Nodes represent a sequence of instructions and edges represent dependences between nodes.

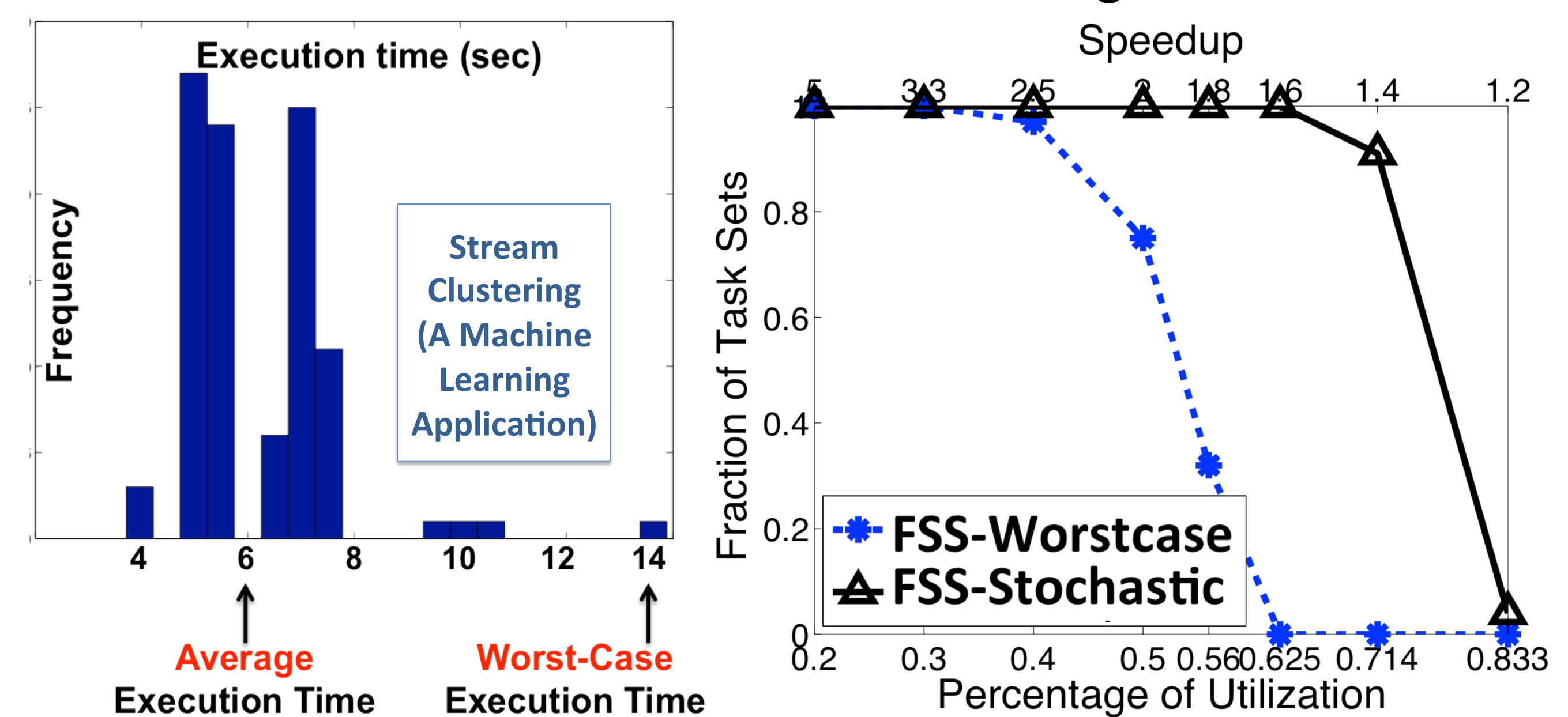
Challenges: Analyze Interference of DAG Tasks

- Dhall's effect: There exist tasks which miss their deadline on m cores even though their utilizations are very close to 1.
- A task can generate jobs with different DAG structures.
- Different DAG structures have different interferences.
- It is computationally expensive to calculate interference of a DAG.



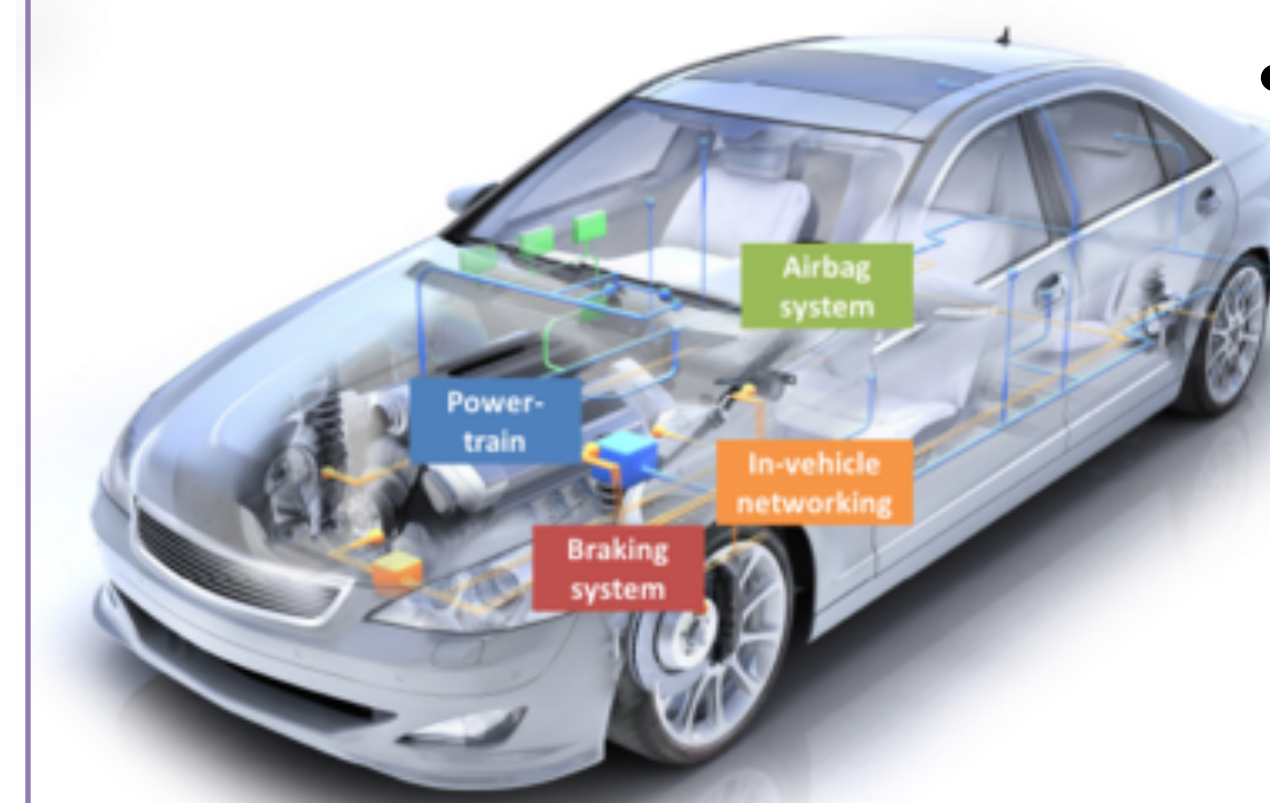
Stochastic Parallel Real-Time Tasks

- If the task parameters have variation from one instance to another, we can model them as random variables and use the mean values for scheduling decisions.



- We generalized federated scheduling for this setting and proved the capacity augmentation bound of 2 [2].

Mixed-Criticality Parallel Tasks



- In autonomous cars, GPS navigation system has lower criticality than braking systems, yet both coexist in one "mixed-criticality" system.
- In normal state, we want to schedule both tasks based on their nominal workload, while in safety-critical state we need to guarantee braking system can still meet its deadline even with the pessimistic worst-case workload.
- Mixed-criticality federated scheduling algorithm has a capacity augmentation bound of 3.732 [6].

Work In Progress and Future Work

- Handling resource contention and critical sections.
- Providing soft real-time guarantees. e.g. tardiness.
- Analyzing and scheduling aperiodic tasks

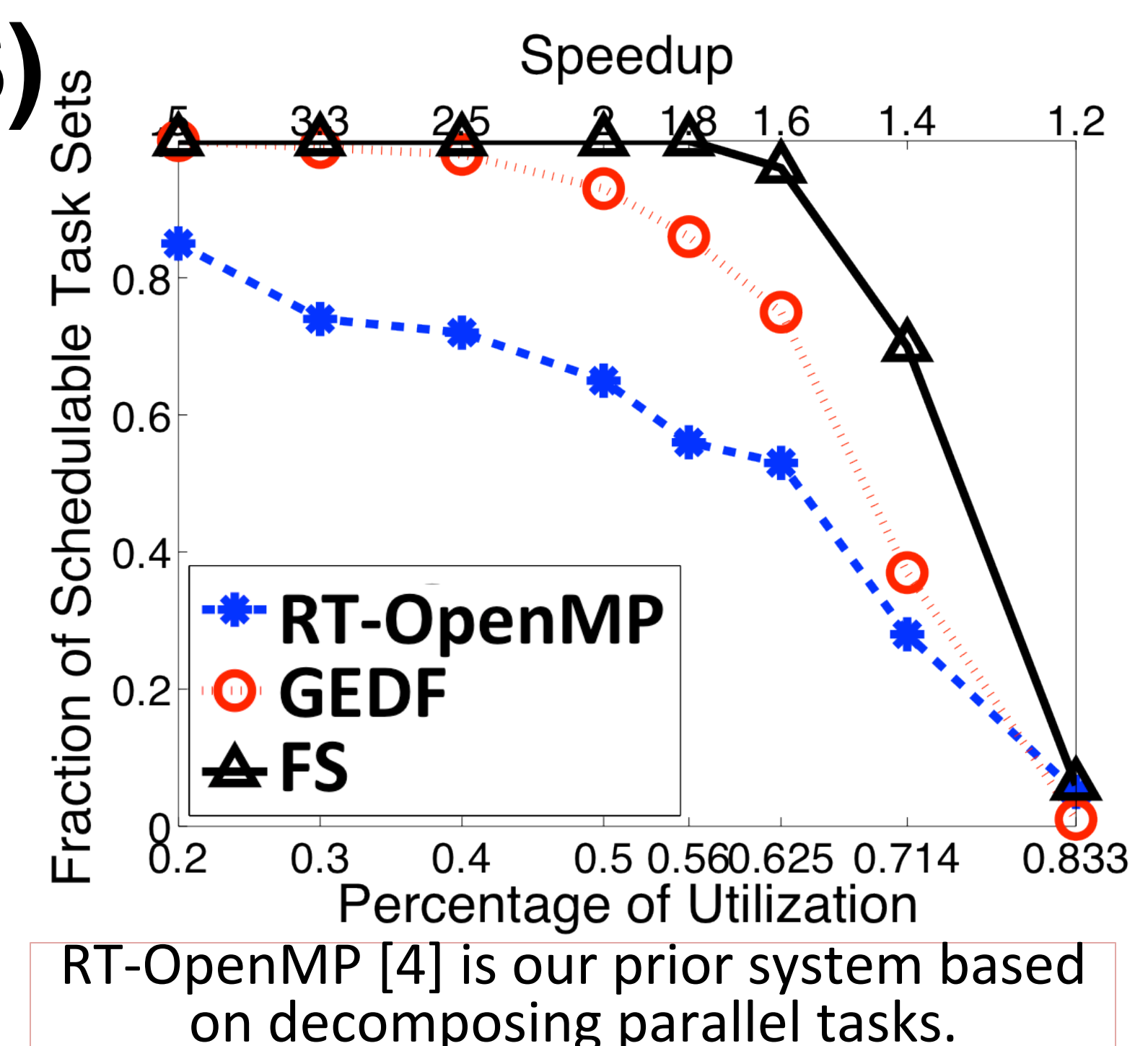
Theoretical Bounds

- Resource augmentation bound of b* : If an ideal scheduler can schedule a task set on m cores of speed 1, then scheduler S can schedule that task set on m cores of speed b

- Capacity augmentation bound of b* : If a task set has total utilization of at most m/b , and each task has its critical-path length of at most $1/b$ of its deadline, then scheduler S that can schedule the task set on m cores.

Global Earliest Deadline First (GEDF) and Federated Scheduling (FS)

	Scheduling Algorithm	Resource Bound	Capacity Bound
GEDF [1, 5]	Global scheduler. A job can execute on all m cores. Prioritizes jobs with the earliest deadline.	approximately 2 (upper & lower bound)	approximately 2.6 (upper & lower bound)
FS [1]	Partition-like $n_i = \left\lceil \frac{C_i - L_i}{D_i - L_i} \right\rceil$	approximately 2	approximately 2 (optimal)



Publications

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- D. Ferry, G. Bunting, A. Megareh, S. Dyk, A. Prakash, K. Agrawal, C. Gill and C. Lu, Real-Time System Support for Hybrid Structural Simulation, EMSOFT, 2014.
- A. Saifullah, D. Ferry, J. Li, K. Agrawal, C. Lu and C. Gill, Parallel Real-Time Scheduling of DAGs, IEEE Transactions

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- J. Li, Z. Luo, D. Ferry, K. Agrawal, C. Gill and C. Lu, Global EDF Scheduling for Parallel Real-Time Tasks, RTS, 2015.
- J. Li, S. Ahuja, K. Agrawal, C. Gill and C. Lu, Mixed-Criticality Federated Scheduling for Real-Time Parallel Tasks (in submission).